

# An introduction to coding theory

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## Lecture #9B: Convolutional codes: Distance properties



# Outline of the lecture

- Weight enumerating function calculation



## Distance properties

- The minimum free distance of a convolutional code

$$d_{\text{free}} \triangleq \min_{\mathbf{u}, \mathbf{u}'} d(\mathbf{v}, \mathbf{v}') : \mathbf{u} \neq \mathbf{u}'$$

where  $\mathbf{v}$  and  $\mathbf{v}'$  are the codewords corresponding to the information sequences  $\mathbf{u}$  and  $\mathbf{u}'$  respectively.



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- $d_{\text{free}}$  is the minimum Hamming distance between any two code sequences in the code.
- $d_{\text{free}}$  is also the minimum weight non-zero sequence, i.e.

$$d_{\text{free}} = \min\{w(\mathbf{v}) : \mathbf{u} \neq 0\}$$



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  - The path gain is the product of the branch gains along a path.
  - The weight of a code sequence is the power of  $X$  in the path gain of the corresponding path.



## Weight enumerating function calculation

- The weight distribution function of a convolutional code is obtained by applying Mason's gain formula to compute the generating function

$$T(X) = \sum_i A_i X^i$$

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- *Forward path*: a path connecting the initial state to the final state that does not go through any state twice.
- *Loop*: a closed path that starts and ends in the same state without going over any other state twice.
- Two or more loops are non-touching if they don't have any states in common.



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- Let  $\{i, j\}$  denote the set of all pairs of non-touching loops.
- Let  $\{i, j, k\}$  denote the set of all triples of non-touching loops, and so on., then define

$$\Delta = 1 - \sum_{\{i\}} C_i + \sum_{\{i, j\}} C_i C_j - \sum_{\{i, j, k\}} C_i C_j C_k + \dots$$





## Weight enumerating function calculation

- Mason's gain formula is given by

$$T(X) = \frac{\sum_i F_i \Delta_i}{\Delta}$$

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- The modified state diagram can be augmented by labeling each branch corresponding to nonzero message bit with Y and labeling every branch with Z.
- The augmented transfer function is given by

$$T(X, Y, Z) = \sum_{i,j,l} A_{i,j,l} X^i Y^j Z^l$$

where  $A_{i,j,l}$  is the number of code sequences of weight  $i$ , whose corresponding information sequence has weight  $j$ , and which has length  $l$  branches.



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- Conditional weight enumerating function (CWEF): Enumerates weights of all codewords associated with particular information weights.



## Weight enumerating function calculation

- For information weight of  $j$ , the CWEF is obtained as

$$T_j(X) = \sum_i A_{i,j} X^i$$

where  $A_{i,j}$  represents number of code sequences with weight  $i$ , and information weight  $j$ .





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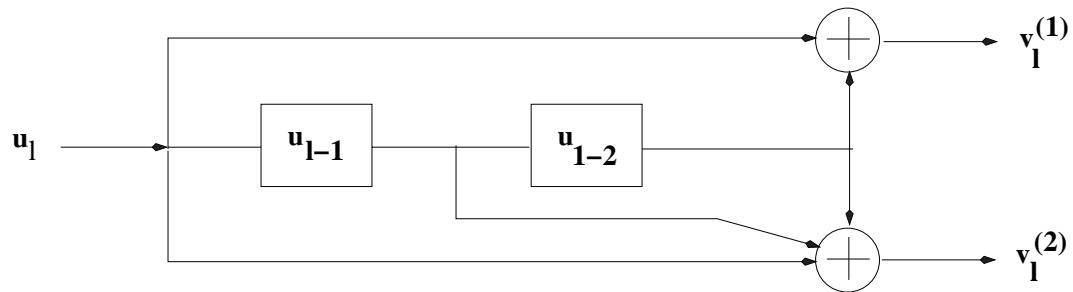
- Input-redundancy weight enumerating function (IRWEF): It is defined for systematic encoders as follows

$$T(W, Y, Z) = \sum_{w,j,l} A_{w,j,l} W^w Y^j Z^l$$

where  $A_{w,j,l}$  is the number of code sequences of length  $l$  with information weight  $j$ , and parity weight  $w$ .

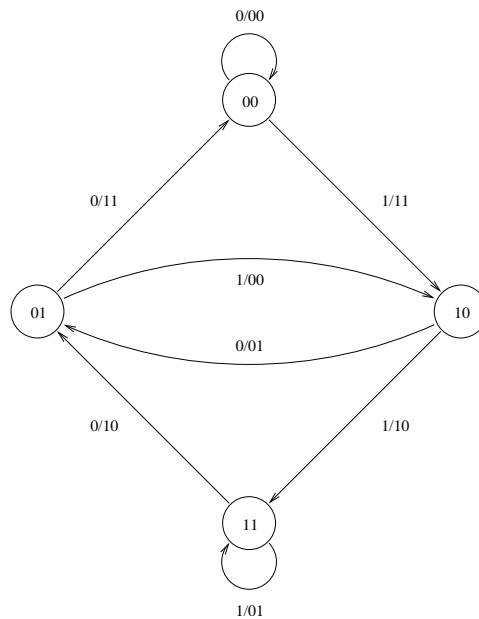


# Convolutional encoder



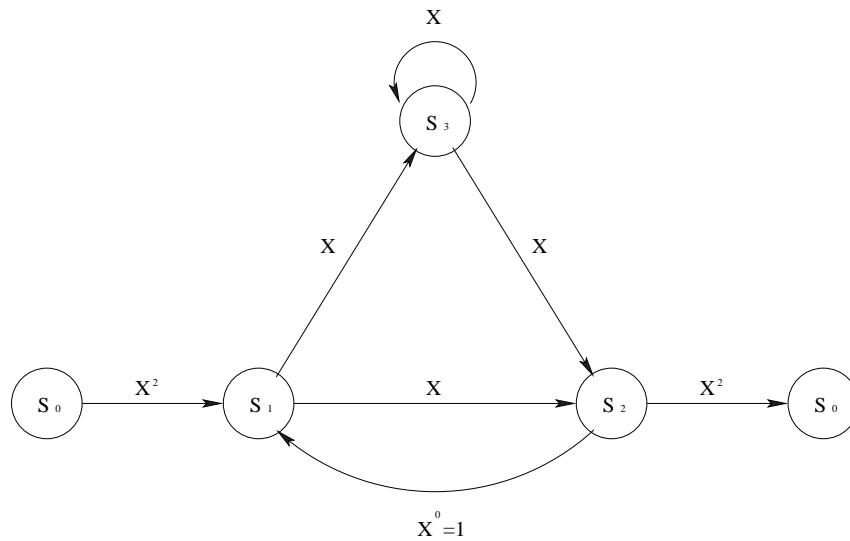
(2, 1, 2) convolutional code

# State diagram



State diagram of (2, 1, 2) convolutional code

## Modified state diagram



Modified state diagram of  $(2, 1, 2)$  convolutional code

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  - $\{1, 2\}$  with  $C_1 C_2 = X^2$ .



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- Pairs of non-touching loops
  - $\{1, 2\}$  with  $C_1 C_2 = X^2$ .
- No triples of non-touching loops.
- Hence,

$$\begin{aligned}\Delta &= 1 - (C_1 + C_2 + C_3) + C_1 C_2 \\ &= 1 - 2X - X^2 + X^2 \\ &= 1 - 2X\end{aligned}$$

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# Weight enumerating function calculation

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## Weight enumerating function calculation

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## Weight enumerating function calculation

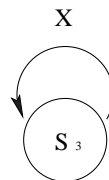
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  - Path 2:  $S_0S_1S_3S_2S_0$  with gain  $F_2 = X^6$ .





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- The graph  $G_1$  obtained by removing the states on the forward path 1 is shown below.



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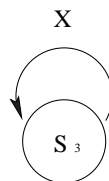
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- $\Delta_1 = 1 - X$ .
- The graph  $G_2$  obtained by removing the states on the forward path 2 is empty.

## Weight enumerating function calculation

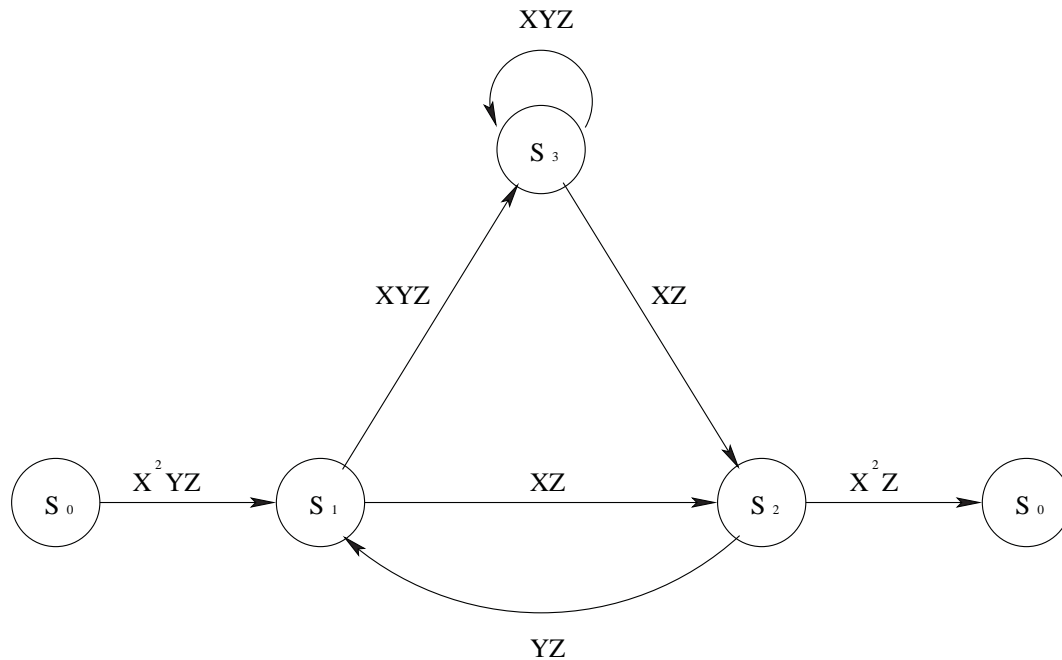
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- $\Delta_1 = 1 - X$ .
- The graph  $G_2$  obtained by removing the states on the forward path 2 is empty.
- Hence  $\Delta_2 = 1$ .



# Augmented state diagram



Navigation icons: back, forward, search, etc.

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# Input-output weight enumerating function calculation

- The transfer function  $T(X, Y, Z)$  is given by

$$\begin{aligned}
 T(X, Y, Z) &= \frac{X^5 Y Z^3}{1 - XYZ - XYZ^2} \\
 &= X^5 Y Z^3 + X^6 (Y^2 Z^4 + Y^2 Z^5) \\
 &\quad + X^7 (Y^3 Z^5 + 2Y^3 Z^6 + Y^3 Z^7) + \dots
 \end{aligned}$$

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