

Numerical Optimization

Constrained Optimization - Algorithms

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NPTEL Course on Numerical Optimization

Barrier and Penalty Methods

Consider the problem:

$$\begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in X \end{array}$$

where $X \in \mathbb{R}^n$.

Idea:

- Approximation by an unconstrained problem
- Solve a sequence of unconstrained optimization problems

Penalty Methods

Penalize for violating a constraint

Barrier Methods

Penalize for reaching the boundary of an inequality constraint

$$\begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in X \end{array}$$

Define a function,

$$\psi(\mathbf{x}) = \begin{cases} 0 & \text{if } \mathbf{x} \in X \\ +\infty & \text{if } \mathbf{x} \notin X \end{cases}$$

Solve an *equivalent* unconstrained problem:

$$\min f(\mathbf{x}) + \psi(\mathbf{x})$$

- *Not a practical approach*
- Replace $\psi(\mathbf{x})$ by a sequence of continuous non-negative functions that approach $\psi(\mathbf{x})$

Penalty Methods

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & \mathbf{x} \in X \end{aligned}$$

- Let \mathbf{x}^* be a local minimum
- Let $X = \{h_j(\mathbf{x}) \leq 0, j = 1, \dots, l\}$
- Define

$$P(\mathbf{x}) = \frac{1}{2} \sum_{j=1}^l [\max(0, h_j(\mathbf{x}))]^2$$

- Define $q(\mathbf{x}, c) = f(\mathbf{x}) + cP(\mathbf{x})$
- Define a sequence $\{c^k\}$ such that $c^k \geq 0$ and $c^{k+1} > c^k \forall k$.
- Let $\mathbf{x}^k = \operatorname{argmin}_{\mathbf{x}} q(\mathbf{x}, c^k)$
- Ideally, $\{\mathbf{x}^k\} \rightarrow \mathbf{x}^*$ as $\{c^k\} \rightarrow +\infty$

Nonlinear Program (NLP)

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & h_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, l \\ & e_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \end{aligned}$$

Define

$$P(\mathbf{x}) = \frac{1}{2} \sum_{j=1}^l [\max(0, h_j(\mathbf{x}))]^2 + \frac{1}{2} \sum_{i=1}^m e_i^2(\mathbf{x})$$

and

$$q(\mathbf{x}, c) = f(\mathbf{x}) + cP(\mathbf{x}).$$

- Assumption: f, h_j 's and e_i 's are *sufficiently smooth*

Lemma

If $\mathbf{x}^k = \operatorname{argmin}_{\mathbf{x}} q(\mathbf{x}, c^k)$ and $c^{k+1} > c^k$, then

- $q(\mathbf{x}^k, c^k) \leq q(\mathbf{x}^{k+1}, c^{k+1})$
- $P(\mathbf{x}^k) \geq P(\mathbf{x}^{k+1})$
- $f(\mathbf{x}^k) \leq f(\mathbf{x}^{k+1})$.

Proof.

$$\begin{aligned}q(\mathbf{x}^{k+1}, c^{k+1}) &= f(\mathbf{x}^{k+1}) + c^{k+1}P(\mathbf{x}^{k+1}) \\ &\geq f(\mathbf{x}^{k+1}) + c^kP(\mathbf{x}^{k+1}) \\ &\geq f(\mathbf{x}^k) + c^kP(\mathbf{x}^k) \\ &= q(\mathbf{x}^k, c^k)\end{aligned}$$

$$\text{Also, } f(\mathbf{x}^k) + c^kP(\mathbf{x}^k) \leq f(\mathbf{x}^{k+1}) + c^kP(\mathbf{x}^{k+1}) \quad \dots (1)$$

$$f(\mathbf{x}^{k+1}) + c^{k+1}P(\mathbf{x}^{k+1}) \leq f(\mathbf{x}^k) + c^{k+1}P(\mathbf{x}^k). \quad \dots (2)$$

Adding (1) and (2), we get $P(\mathbf{x}^k) \geq P(\mathbf{x}^{k+1})$.

$$f(\mathbf{x}^{k+1}) + c^kP(\mathbf{x}^{k+1}) \geq f(\mathbf{x}^k) + c^kP(\mathbf{x}^k) \Rightarrow f(\mathbf{x}^{k+1}) \geq f(\mathbf{x}^k) \quad \square$$

Lemma

Let \mathbf{x}^* be a solution to the problem,

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & \mathbf{x} \in X. \end{aligned} \quad \dots (P1)$$

Then, for each k , $f(\mathbf{x}^k) \leq f(\mathbf{x}^*)$.

Proof.

$$\begin{aligned} f(\mathbf{x}^k) &\leq f(\mathbf{x}^k) + c^k P(\mathbf{x}^k) \\ &\leq f(\mathbf{x}^*) + c^k P(\mathbf{x}^*) = f(\mathbf{x}^*) \end{aligned}$$



Theorem

Any limit point of the sequence, $\{\mathbf{x}^k\}$ generated by the penalty method is a solution to the problem (P1).

Nonlinear Program (NLP)

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & h_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, l \\ & e_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \end{aligned}$$

Penalty Function Method (to solve NLP)

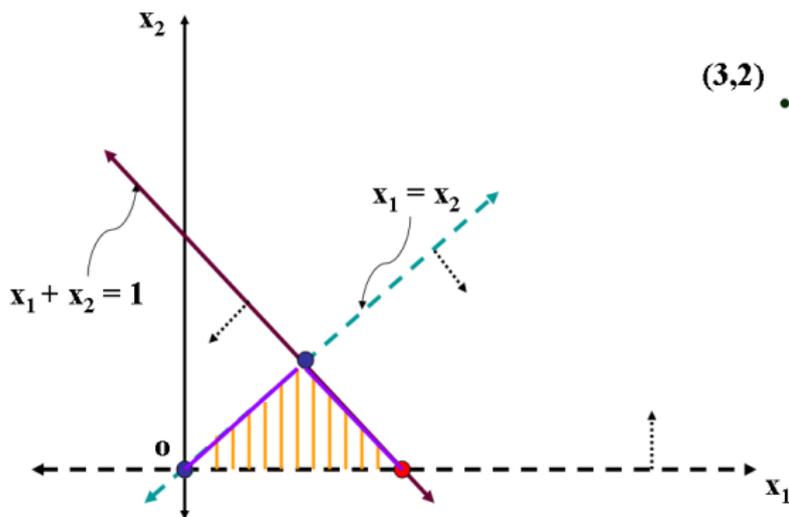
- (1) Input: $\{c^k\}_{k=0}^{\infty}, \epsilon$
 - (2) Set $k := 0$, initialize \mathbf{x}^k
 - (3) **while** $(q(\mathbf{x}^k, c^k) - f(\mathbf{x}^k)) > \epsilon$
 - (a) $\mathbf{x}^{k+1} = \operatorname{argmin}_{\mathbf{x}} q(\mathbf{x}, c^k)$
 - (b) $k := k + 1$
- endwhile**

Output : $\mathbf{x}^* = \mathbf{x}^k$

Example:

$$\begin{aligned} \min \quad & \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] \\ \text{s.t.} \quad & -x_1 + x_2 \leq 0 \\ & x_1 + x_2 \leq 1 \\ & -x_2 \leq 0 \end{aligned}$$

$$\mathbf{x}^* = (1, 0)^T$$



$$\begin{aligned}
 \min \quad & \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] \\
 \text{s.t.} \quad & -x_1 + x_2 \leq 0 \\
 & x_1 + x_2 \leq 1 \\
 & -x_2 \leq 0
 \end{aligned}$$

$$\begin{aligned}
 q(\mathbf{x}, c) = \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] + \frac{c}{2} & [(\max(0, -x_1 + x_2))^2 \\
 & + (\max(0, x_1 + x_2 - 1))^2 + (\max(0, -x_2))^2]
 \end{aligned}$$

- Let $\mathbf{x}^0 = (3, 2)^T$ (Violates the constraint $x_1 + x_2 \leq 1$)
- At \mathbf{x}^0 ,

$$q(\mathbf{x}, c) = \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] + \frac{c}{2}[(x_1 + x_2 - 1)^2].$$

- $\nabla_x q(\mathbf{x}, c) = \mathbf{0} \Rightarrow \mathbf{x}^1(c) = \begin{pmatrix} \frac{2c+3}{2} \\ \frac{2c+1}{2c+1} \end{pmatrix} = \mathbf{x}^*(c)$

$$\begin{aligned}
 \min \quad & \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] \\
 \text{s.t.} \quad & -x_1 + x_2 \leq 0 \\
 & x_1 + x_2 \leq 1 \\
 & -x_2 \leq 0
 \end{aligned}$$

- At $\mathbf{x}^0 = (3, 2)^T$,

$$q(\mathbf{x}, c) = \frac{1}{2}[(x_1 - 3)^2 + (x_2 - 2)^2] + \frac{c}{2}[(x_1 + x_2 - 1)^2].$$

- $\nabla_x q(\mathbf{x}, c) = \mathbf{0} \Rightarrow \mathbf{x}^1(c) = \begin{pmatrix} \frac{2c+3}{2c+1} \\ \frac{2}{2c+1} \end{pmatrix} = \mathbf{x}^*(c)$
- Taking limit as $c \rightarrow \infty$, $\mathbf{x}^* = (1, 0)^T$

Consider the problem,

$$\min \quad f(\mathbf{x})$$

- Let (\mathbf{x}^*, μ^*) be a KKT point (s.t. $e(\mathbf{x}) \equiv 0$, $\nabla f(\mathbf{x}^*) + \mu^* \nabla e(\mathbf{x}^*) = \mathbf{0}$)
- Penalty Function: $q(\mathbf{x}, c) = f(\mathbf{x}) + cP(\mathbf{x})$
- As $c \rightarrow \infty$, $q(\mathbf{x}, c) = f(\mathbf{x})$

Consider the *perturbed* problem,

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & e(\mathbf{x}) = \theta \end{aligned}$$

and the penalty function,

$$\begin{aligned} \hat{q}(\mathbf{x}, c) &= f(\mathbf{x}) + c(e(\mathbf{x}) - \theta)^2 \\ &= f(\mathbf{x}) - 2c\theta e(\mathbf{x}) + ce(\mathbf{x})^2 \quad (\text{ignoring constant term}) \\ &= \underbrace{f(\mathbf{x}) + \mu e(\mathbf{x})}_{\mathcal{L}(x, \mu)} + ce(\mathbf{x})^2 \\ &= \hat{\mathcal{L}}(\mathbf{x}, \mu, c) \quad (\text{Augmented Lagrangian Function}) \end{aligned}$$

At (\mathbf{x}^*, μ^*) , $\nabla_x \mathcal{L}(\mathbf{x}^*, \mu^*) = \nabla f(\mathbf{x}^*) + \mu^* \nabla e(\mathbf{x}^*) = \mathbf{0}$.

$$\begin{aligned} \therefore \nabla_x \hat{q}(\mathbf{x}^*, c) &= \nabla_x \hat{\mathcal{L}}(\mathbf{x}^*, \mu^*, c) \\ &= \nabla_x \mathcal{L}(\mathbf{x}^*, \mu^*) + 2ce(\mathbf{x}^*) \nabla e(\mathbf{x}^*) \\ &= \mathbf{0} \quad \forall c \end{aligned}$$

Q. How to get an estimate of μ^* ?

Let \mathbf{x}_c^* be a minimizer of $\mathcal{L}(\mathbf{x}, \mu, c)$. Therefore,

$$\begin{aligned} \nabla_x \mathcal{L}(\mathbf{x}_c^*, \mu, c) &= \nabla f(\mathbf{x}_c^*) + \mu \nabla e(\mathbf{x}_c^*) + ce(\mathbf{x}_c^*) \nabla e(\mathbf{x}_c^*) = \mathbf{0} \\ \therefore \nabla f(\mathbf{x}_c^*) &= - \underbrace{(\mu + ce(\mathbf{x}_c^*))}_{\text{estimate of } \mu^*} \nabla e(\mathbf{x}_c^*) \end{aligned}$$

Program (EP)

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & e(\mathbf{x}) = 0 \end{aligned}$$

Augmented Lagrangian Method (to solve EP)

- (1) Input: c, ϵ
- (2) Set $k := 0$, initialize \mathbf{x}^k, μ^k
- (3) **while** $(\hat{\mathcal{L}}(\mathbf{x}^k, \mu^k, c) - f(\mathbf{x}^k)) > \epsilon$
 - (a) $\mathbf{x}^{k+1} = \operatorname{argmin}_{\mathbf{x}} \hat{\mathcal{L}}(\mathbf{x}, \mu^k, c)$
 - (b) $\mu^{k+1} = \mu^k + c e(\mathbf{x}^k)$
 - (c) $k := k + 1$

endwhile

Output : $\mathbf{x}^* = \mathbf{x}^k$

Nonlinear Program (NLP)

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & h_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, l \\ & e_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \end{aligned}$$

- Easy to extend the Augmented Lagrangian Method to NLP
- Rewrite the inequality constraint, $h(\mathbf{x}) \leq 0$ as an equality constraint,

$$h(\mathbf{x}) + y^2 = 0$$

Barrier Methods

- Typically applicable to inequality constrained problems

$$\begin{array}{ll} \min & f(\mathbf{x}) \\ \text{s.t.} & h_j(\mathbf{x}) \leq 0, j = 1, \dots, l \end{array}$$

Let $X = \{\mathbf{x} : h_j(\mathbf{x}) \leq 0, j = 1, \dots, l\}$

- Some Barrier functions (defined on the *interior* of X)

$$B(\mathbf{x}) = -\sum_{j=1}^l \frac{1}{h_j(\mathbf{x})} \quad \text{or} \quad B(\mathbf{x}) = -\sum_{j=1}^l \log(-h_j(\mathbf{x}))$$

- Approximate problem using Barrier function (for $c > 0$)

$$\begin{array}{ll} \min & f(\mathbf{x}) + \frac{1}{c}B(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in \text{Interior of } X \end{array}$$

Cutting-Plane Methods

Primal Problem

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & h_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, l \\ & e_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \\ & \mathbf{x} \in X \end{aligned}$$

X is a compact set.

Dual Function: $z = \theta(\boldsymbol{\lambda}, \boldsymbol{\mu}) = \min_{\mathbf{x} \in X} f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x})$

Dual Problem

$$\begin{aligned} \max \quad & \theta(\boldsymbol{\lambda}, \boldsymbol{\mu}) \\ \text{s.t.} \quad & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

Equivalent Dual problem

$$\begin{aligned} \max_{z, \boldsymbol{\mu}, \boldsymbol{\lambda}} \quad & z \\ \text{s.t.} \quad & z \leq f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x}), \quad \mathbf{x} \in X \\ & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

Linear Program with infinite constraints

Equivalent Dual problem

$$\begin{aligned} \max_{z, \mu, \lambda} \quad & z \\ \text{s.t.} \quad & z \leq f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x}), \quad \mathbf{x} \in X \\ & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

Idea: Solve an approximate dual problem.

Suppose we know $\{\mathbf{x}^j\}_{j=0}^{k-1}$ such that

$$z \leq f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x}), \quad \mathbf{x} \in \{\mathbf{x}^0, \dots, \mathbf{x}^{k-1}\}$$

Approximate Dual Problem

$$\begin{aligned} \max_{z, \mu, \lambda} \quad & z \\ \text{s.t.} \quad & z \leq f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x}), \quad \mathbf{x} \in \{\mathbf{x}^0, \dots, \mathbf{x}^{k-1}\} \\ & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

Let $(z^k, \boldsymbol{\lambda}^k, \boldsymbol{\mu}^k)$ be the optimal solution to this problem.

Approximate Dual Problem

$$\begin{aligned} \max_{z, \mu, \lambda} \quad & z \\ \text{s.t.} \quad & z \leq f(\mathbf{x}) + \boldsymbol{\lambda}^T h(\mathbf{x}) + \boldsymbol{\mu}^T e(\mathbf{x}), \quad \mathbf{x} \in \{\mathbf{x}^0, \dots, \mathbf{x}^{k-1}\} \\ & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

If $z^k \leq f(\mathbf{x}) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}) + \boldsymbol{\mu}^{kT} e(\mathbf{x}) \forall \mathbf{x} \in X$, then $(z^k, \boldsymbol{\lambda}^k, \boldsymbol{\mu}^k)$ is the solution to the dual problem.

Q. How to check if $z^k \leq f(\mathbf{x}) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}) + \boldsymbol{\mu}^{kT} e(\mathbf{x}) \forall \mathbf{x} \in X$?
Consider the problem,

$$\begin{aligned} \min \quad & f(\mathbf{x}) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}) + \boldsymbol{\mu}^{kT} e(\mathbf{x}) \\ \text{s.t.} \quad & \mathbf{x} \in X \end{aligned}$$

and let \mathbf{x}^k be an optimal solution to this problem.

$$\begin{array}{ll} \min & f(\mathbf{x}) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}) + \boldsymbol{\mu}^{kT} e(\mathbf{x}) \\ \text{s.t.} & \mathbf{x} \in X \end{array}$$

and let \mathbf{x}^k be an optimal solution to this problem.

- If $z^k \leq f(\mathbf{x}^k) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}^k) + \boldsymbol{\mu}^{kT} e(\mathbf{x}^k)$, then $(\boldsymbol{\lambda}^k, \boldsymbol{\mu}^k)$ is an optimal solution to the Lagrangian dual problem.
- If $z^k > f(\mathbf{x}^k) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}^k) + \boldsymbol{\mu}^{kT} e(\mathbf{x}^k)$, then add the constraint, $z \leq f(\mathbf{x}^k) + \boldsymbol{\lambda}^T h(\mathbf{x}^k) + \boldsymbol{\mu}^T e(\mathbf{x}^k)$ to the approximate dual problem.

Nonlinear Program (NLP)

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & h_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, l \\ & e_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \\ & \mathbf{x} \in X \end{aligned}$$

Summary of steps for Cutting-Plane Method:

- Initialize with a feasible point \mathbf{x}^0
- while stopping condition is not satisfied

$$\begin{aligned} (z^k, \boldsymbol{\lambda}^k, \boldsymbol{\mu}^k) = \operatorname{argmax}_{z, \boldsymbol{\lambda}, \boldsymbol{\mu}} \quad & z \\ \text{s.t.} \quad & z \leq f(\mathbf{x}^j) + \boldsymbol{\lambda}^T h(\mathbf{x}^j) + \boldsymbol{\mu}^T e(\mathbf{x}^j), \quad j = 0, \dots, k-1 \\ & \boldsymbol{\lambda} \geq \mathbf{0} \end{aligned}$$

$$\mathbf{x}^k = \operatorname{argmin}_{\mathbf{x} \in X} f(\mathbf{x}) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}) + \boldsymbol{\mu}^{kT} e(\mathbf{x})$$

Stop if $z^k \leq f(\mathbf{x}^k) + \boldsymbol{\lambda}^{kT} h(\mathbf{x}^k) + \boldsymbol{\mu}^{kT} e(\mathbf{x}^k)$. Else, $k := k + 1$.