

The Lecture Contains:

- ☰ Semaphore Types
- ☰ Synchronization Among Processes
- ☰ OS Implementations
- ☰ Linux API for Semaphores
- ☰ Pthread Mutex Locks
- ☰ Win32 Mutex Locks and Semaphores
- ☰ Processes in a Group
- ☰ Cooperating Processes (example)
- ☰ Cooperating Processes
- ☰ Shared Memory System
- ☰ Message Passing (Send)
- ☰ Message Passing (Receive)
- ☰ Cooperating Processes
- ☰ IPC: Shared Memory
- ☰ Producer and Consumer Code

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## Module 21: Problem and Solution

## Lecture 42: Multi-core ComputingInter-process Communication

## Semaphore Types

- Binary or counting
- Binary: Semaphore value can be T or F only
- Binary semaphores provide mutual exclusion
  - Sometimes known as mutex locks
- Counting semaphores
  - Can be used when the given resource has finite number of instances (n).
  - Initialize semaphore.value to n.

## Synchronization Among Processes

- Two Processes P1 and P2.
- How do we make sure that P1 executes S1 first before P2 executes S2?

semaphore synch;

P1:

S1;

synch.signal();

P2:

synch.wait();

S2

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## Module 21: Problem and Solution

## Lecture 42: Multi-core Computing Inter-process Communication

## OS Implementations

- Busy wait is not tolerated in an OS!!
- Bounded wait is to be ensured.
- Solution:
  - Use sleep and wakeup to move processes from running to waiting state and waiting to ready state.
  - Maintain a queue of processes waiting and wake up processes in that order.

```
class semaphore {
private:
int value;
list<PCB> wait_list;
public:
semaphore() {
value = 0;
wait_list = list<PCB>();
}
semaphore(int a) {
value = a;
wait_list = list<PCB>();
}
```

```
void wait(void) {
value--;
if (value < 0) { wait_list.push_front(
current);
sleep();
}
}
void signal(void) {
PCB p;
value++;
if (value <= 0) {
p = *(wait_list.end()); wait_list.pop_back();
wakeup(p);
}
}
}
```

Wait also known as *P*  
Signal also known as *V*

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## Module 21: Problem and Solution

## Lecture 42: Multi-core Computing Inter-process Communication

## Linux API for Semaphores

- Semget
  - To get an array of semaphores
- Semctl
  - Semaphore controls. For example initial value
- Semop
  - Semaphore Operations (lock, signal etc.)

## Pthread Mutex Locks

- Data Types: pthread\_mutex\_t
  - Creation: pthread\_mutex\_init
  - Lock: pthread\_mutex\_lock
  - Unlock: pthread\_mutex\_unlock
- ```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
pthread_mutex_lock(&mutex);
pthread_mutex_unlock(&mutex);
```

## Win32 Mutex Locks and Semaphores

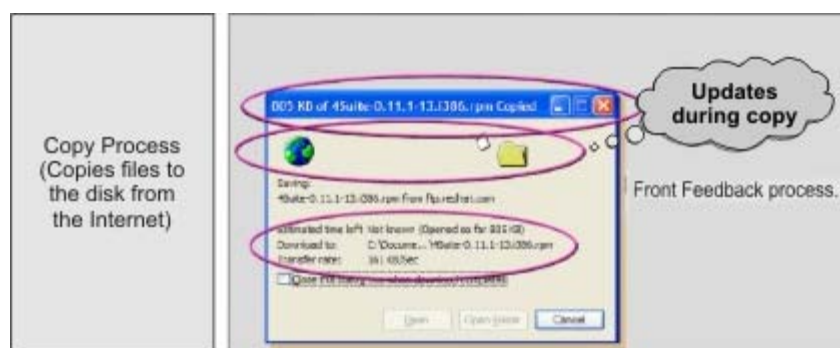
- Declaration
  - HANDLE mutex; (or HANDLE semaphore);
- Creation
  - CreateMutex: To create a mutex lock
  - CreateSemaphore: To create a semaphore
    - semaphore = CreateSemaphore(NULL, 1, 5, NULL);
- Wait for mutex or semaphore
  - WaitForSingleObject(HANDLE, WAITTYPE)
    - WaitForSingleObject(Sem, INFINITE);
- Releasing
  - ReleaseMutex(mutex)
  - ReleaseSemaphore(sem, 1, NULL)

## Multi-core ComputingInter-process Communication

### Processes in a Group

- A process can be independent
  - Is not directly affected by other processes.
  - Does not affect other processes.
  - Example: /bin/lis and the shell
    - Are they related?
- Processes may be cooperating
  - Information Sharing
  - Speed up of execution
  - Modularity and convenience

### Cooperating Processes (example)

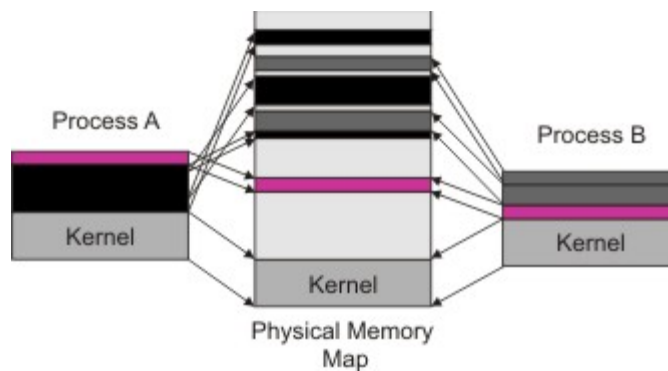


- Not really an example of “processes” but “threads”.
- The issues are the similar though.

## Cooperating Processes

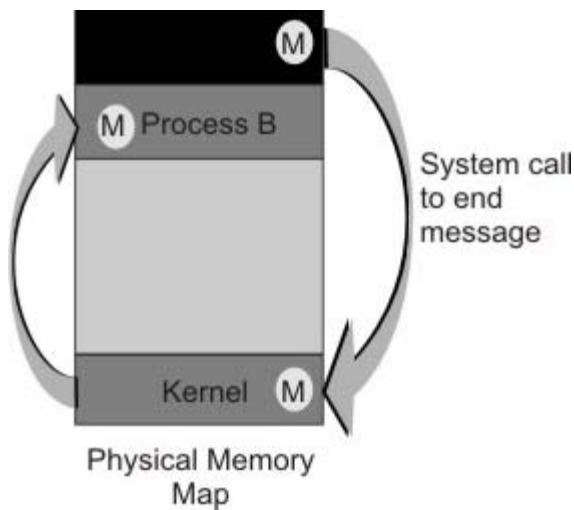
- Require
  - Inter process communication
    - Shared memory between processes
    - Message passing
      - Sender makes a call to the OS to send a message
      - Receiver makes a call to read message from the OS
  - Producer consumer relationship
    - A process produces data to be consumed by other process.

## Shared Memory System



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## Message Passing (Send)



## Message Passing (Receive)

- Receive can be blocking
  - A process makes a system call to receive a message.
    - If message is not available, the process is made to sleep (wait) and woken up when message is received.
- Receive can be non-blocking
  - Process makes a system call to receive a message.
    - Return value from the system call determines whether a message is ready or not.

## Module 21: Problem and Solution

## Lecture 42: Multi-core ComputingInter-process Communication

## Cooperating Processes

- Inter process communication
  - Shared memory between processes
  - Message passing
- Producer consumer relationship

## IPC: Shared Memory

- Shared buffer between processes
- ```
#define BUF_SZ 1024
typedef struct {
    ...
} BUF_Data;
struct {
    BUF_Data items[BUF_SZ];
    int inptr, outptr; /* Global variables */
} buffer; /* Must be shared between
/* two processes */
```

## Producer and Consumer Code

```
void produce(BUF_Data item) {
    while ((buffer.inptr+1)%BUF_SZ == buffer.outptr) ;
    buffer.items[buffer.inptr] = item;
    buffer.inptr = (buffer.inptr +1)%BUF_SZ;
}
BUF_Data consume(void) {
    BUF_Data item;
    while (buffer.outptr == buffer.inptr) ;
    item = buffer.items[buffer.outptr];
    buffer.outptr = (buffer.outptr +1)%BUF_SZ;
    return (item);
}
```

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