

The Lecture Contains:

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- Scheduling
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- SJF
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- Preemptive SJF
- Example (SRTF)
- Adding Priority (Priority Scheduling)
- Priority Scheduling + SJF

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Multi-core computing Operating Systems

Contents

- Scheduling
 - Traditional
 - Real time
 - Multi-processor
- Synchronization and synchronization objects
- Inter-process communication
- Security

Multi-core computing CPU Scheduling

Scheduling

- To hide the effects of I/O Bursts and achieve higher CPU utilization
- To give a fair chance to all processes
- Short term scheduling
 - CPU scheduling
- Long term scheduling
 - Process admission policies
- Medium term scheduling
 - Swap management

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Module 19: Multi-core computing Operating Systems

Lecture 37: Multi-core computing CPU Scheduling

CPU Scheduling

- CPU scheduler selects a process from “Ready to run” queue
 - Scheduling decision
- Allocates CPU to this process.
 - Dispatch
 - Context switching
- Queue is maintained by PCB pointers
- Pre-emptive scheduling
 - When a running process may be removed and put back in the “ready to run” queue.

When to Have CPU Scheduling

- Process moves from running state to waiting state
 - Due to an I/O request
 - Due to a call to wait
- Upon Process Termination
- Expiration of time quota (preemptive)
- Upon any other time when OS is called
 - Interrupts, System Calls

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Module 19: Multi-core computing Operating Systems

Lecture 37: Multi-core computing CPU Scheduling

Scheduling Criteria

- CPU Utilization (System centric)
 - Keep CPU as much busy as possible
- Throughput (System centric)
 - Number of processes completed/time
- Turnaround time (Process centric)
 - Real time taken to complete a process
- Waiting time (Process centric)
 - How much time a process is in ready queue
- Response time (Process centric)
 - Factor for an interactive process
- Deadline (Real time behavior)
 - Time guarantee to schedule a task.

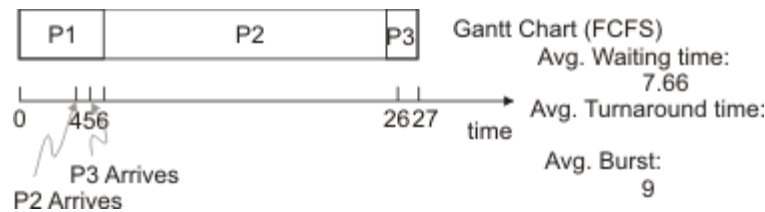
Scheduling Algorithms

- First Come First Serve (FCFS)
 - Simplest implementation
 - No preemption
- Shortest Job First (SJF)
 - Optimal scheduling algorithm
 - Minimum Average Waiting Time
 - Difficulty: To know the time that it will take
 - Batch systems: A good choice.

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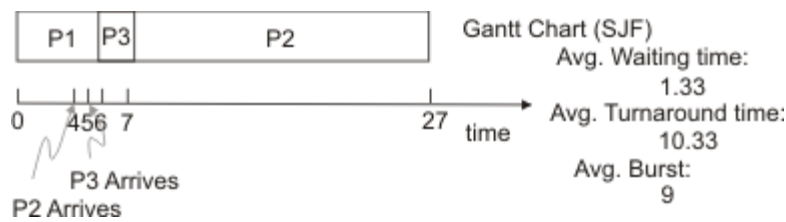
Example

Process	Arrival Time	CPU Burst	Wait Time	Turnaround
P1	0	6	0	6
P2	4	20	2	22
P3	5	1	21	22



Example (Contd.)

Process	Arrival Time	CPU Burst	Wait Time	Turnaround
P1	0	6	0	6
P2	4	20	3	23
P3	5	1	1	2



SJF

- Though optimal, not practical.
 - No way to find the CPU burst.
- Scheduling techniques try to approximate SJF
 - Guess/predict the next burst (τ_n : Predicted n^{th} burst. t_n : Actual n^{th} burst)

$$t_{n+1} = \alpha t_n + (1 - \alpha)\tau_n \quad 0 \leq \alpha \leq 1$$
$$\tau_{n+1} = \alpha t_n + (1 - \alpha)\tau_n$$
$$\tau_{n+1} = t_n$$
$$\tau_{n+1} = \tau_n$$
$$\tau_{n+1} = \alpha t_n + (1 - \alpha) \cdot \alpha \cdot t_{n-1} + \dots + (1 - \alpha)^j \cdot \alpha \cdot t_{n-j} + \dots + (1 - \alpha)^{n+1} \tau_0$$

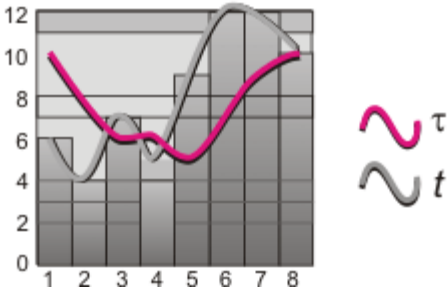
$$0 \leq \alpha \leq 1$$

when $\alpha = 1$
when $\alpha = 0$

- Burst guess is obtained by exponential averaging.
- τ_0 is a system wide scheduling constant.
- α is a weight factor.

Example of “Guessing”

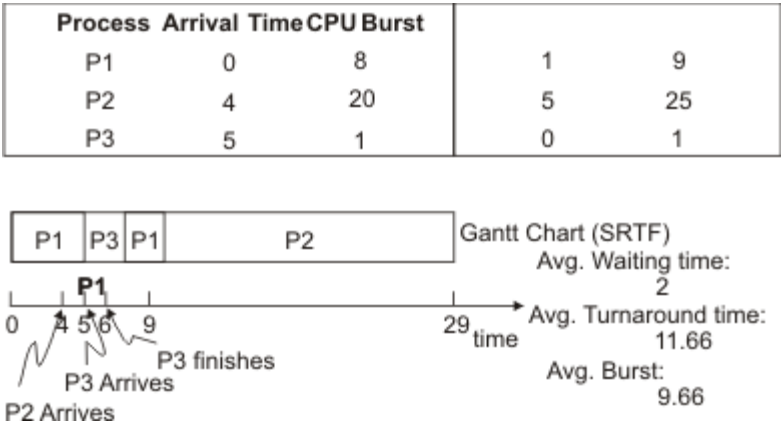
CPU Burst	6	4	7	5	9	12	12	10
Guess ($\tau_0=10$, $\alpha=0.5$)	10	8	6	6	9	7	9	10



Preemptive SJF

- When a new task arrives, the SJF is evaluated again and rescheduling can take place.
- aka Shortest-remaining-time-first (SRTF) scheduling

Example (SRTF)



Adding Priority (Priority Scheduling)

- Simplest form is to add priority to a process.
 - Select the highest priority job first
 - Apply SJF/SRTF/FIFO/... scheduling only within a group of processes with the same priority

Priority scheduling + SJF

