

Module 5: Disk-based Index Structures

Lecture 21: R-tree Variants

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The Lecture Contains:

☰ Greene's R-tree

☰ R*-tree

- Insertion
- Splitting
- Forced reinsertion

☰ R + - tree

- Example
- Searching

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Greene's R-tree

- Different splitting strategy than R-tree
- Choose axis to split
 - Find two most distant rectangles (seeds)
 - For each axis, record separation of seeds
 - Normalize separation by range of axis
 - Choose axis with largest separation
- Distribute regions
 - Sort entries by min value along chosen axis
 - Assign first $(M + 1)/2$ entries to one group, last $(M + 1)/2$ entries to another group
 - Assign remaining entry (if $M + 1$ is odd) to group with largest increase in volume

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R*-tree

- Three optimizations over R-trees
 - Insertion
 - Splitting
 - Forced reinsertion when over flow
- Instead of just volume, four parameters:
 - Volume wasted in a node ("dead space")
 - Volume of overlaps
 - Margin (i.e., sum of edges)
 - Space utilization

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Insertion

- If child pointers are not leaf nodes
 - Choose subtree with least volume increase (as before)
- If child pointers are leaf nodes (level leaf-1)
 - Choose subtree with least overlap enlargement
 - Overlap (E_k) = sum of volume($E_k \setminus E_i$) for all other entries E_i
 - $O(M^2)$ because $O(M)$ entries need to be checked and each entry, overlap to be calculated for $O(M)$ siblings

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Splitting

- Two parameters:
 - Margin: $\text{Margin}(\text{MBR}(\text{Group 1})) + \text{Margin}(\text{MBR}(\text{Group 2}))$
 - Overlap: $\text{Volume}(\text{MBR}(\text{Group 1}) \cap \text{MBR}(\text{Group 2}))$
- Choose split axis
 - For each axis, sort entries by low values
 - Enumerate $M - 2m + 2$ possible groupings or distributions

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Splitting

- Two parameters:
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 - Overlap: $\text{Volume}(\text{MBR}(\text{Group } 1) \cap \text{MBR}(\text{Group } 2))$
- Choose split axis
 - For each axis, sort entries by low values
 - Enumerate $M - 2m + 2$ possible groupings or distributions
 - k^{th} distribution: First $m - 1 + k$ entries and rest
 - Compute sum of margins along that axis of all distributions
 - Do the same by sorting high values of each axis
 - Choose axis with least sum of margin values
 - $O(2d.M \lg M)$
- For chosen axis and low/high value, choose distribution with least overlap

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Forced reinsertion

- When over flow occurs, do not split immediately
- Choose p entries to be re-inserted
 - Empirically, $p = 30\%$
- Which p entries?
- For each entry, compute distance of its centroid to centroid of node
 - Remove from closest to farthest (close reinsert)
 - Reinsertion will not increase volume of node
 - Remove from farthest to closest (far reinsert)
 - Likely to be reinserted at some other node
 - Not clear which one better
 - Experimentally, close reinsert better
- $O(N \lg N)$
- Done only once at each level to avoid infinite loop

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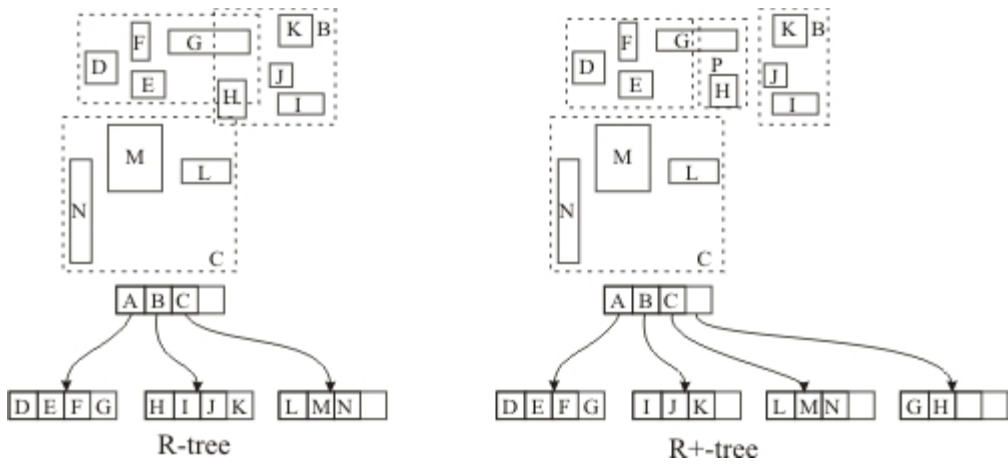
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R + - tree

- Mix of R-tree's data partitioning with K-d-B-tree's space-partitioning
- Reduces dead space problem of K-d-B trees by using index rectangles
- Reduces node overlap by making siblings disjoint
- Data objects split into rectangles
 - Disjoint
 - Union gives original data object
- Data object stored in multiple leaves
- Downward splitting
- No space utilization guarantee

Example

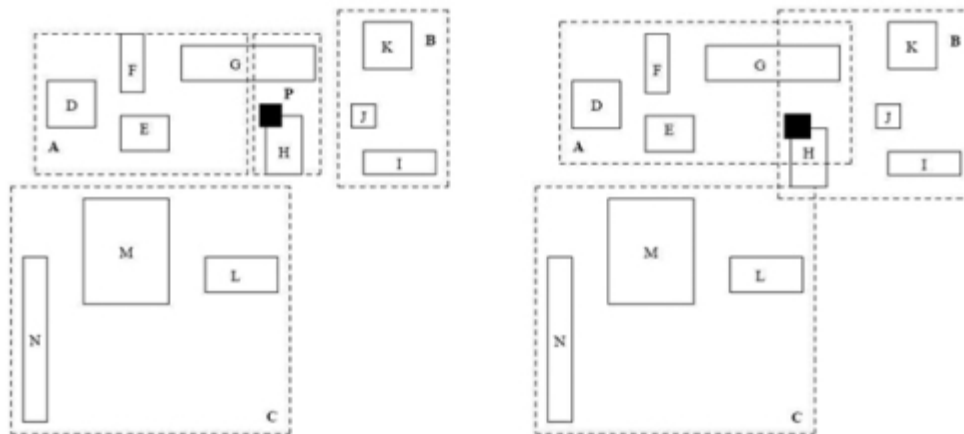


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Searching



- In R-tree, nodes A and B are searched
- In R+-tree, only node P is searched
- Trade-off between overlapped searches and increased height
- Not much better than R*-trees