

Unit 8 - Week 7: Properties of measures of Information-2 and Information Theoretic lower bounds

Course outline
How does an NPTEL online course work?
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Week 2: Uncertainty, compression, and entropy
Week 3: Randomness and entropy
Week 4: Information and statistical inference-1
Week 5: Information and statistical inference-2
Week 6: Properties of measures of Information-1
Week 7: Properties of measures of Information-2 and Information Theoretic lower bounds
<ul style="list-style-type: none"> ● Midyear Review ● Proof of Fano's inequality <input type="radio"/> Variational formulae <input type="radio"/> Capacity as information radius <input type="radio"/> Proof of Pinsker's inequality <input type="radio"/> Continuity of entropy ● Lower bound for compression <input type="radio"/> Lower bound for hypothesis testing <input type="radio"/> Quiz : Weekly Assignment 7 <input type="radio"/> Unit-7 notes <input type="radio"/> Solution 7
Week 8: Information theoretic lower bounds and Data Compression-1
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Week 12: Shannon's channel coding theorem proof (Unit 13) and Gaussian channels (Unit 14)
VIDEO DOWNLOAD
Live Session

Weekly Assignment 7

The due date for submitting this assignment has passed.
As per our records you have not submitted this assignment.

Due on 2020-11-04, 23:59 IST.

Unless stated otherwise, assume that probability distributions are defined over a finite alphabet, and that 'log' is to the base 2.

1) Let X, Y be random variables with joint distribution P_{XY} , where X is distributed uniformly over $\{1, \dots, M\}$. Let $g: \mathcal{Y} \rightarrow \mathcal{X}$ be a deterministic function **1 point** and let $p_e = P_{XY}(X \neq g(Y))$. Which of the following are true? ($\text{Ber}(p)$ denotes the Bernoulli distribution with parameter p .)

- $I(X \wedge Y) \geq D(\text{Ber}(1 - p_e) \parallel \text{Ber}(\frac{1}{M}))$
- $I(X \wedge Y) \leq D(\text{Ber}(p_e) \parallel \text{Ber}(1 - \frac{1}{M}))$
- If $1 - p_e > \frac{1}{M}$, then X and Y are not independent.
- If X and Y are not independent, then $1 - p_e > \frac{1}{M}$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$I(X \wedge Y) \geq D(\text{Ber}(1 - p_e) \parallel \text{Ber}(\frac{1}{M}))$$

If $1 - p_e > \frac{1}{M}$, then X and Y are not independent.

2) Let $\mathcal{X} = \{1, 2, 3\}$. Let $W: \mathcal{X} \rightarrow \mathcal{X}$ be a channel defined as follows: $W_1 = (0, 1, 0)$, $W_2 = (0, 0, 1)$, $W_3 = (1, 0, 0)$. Let P and Q be two **1 point** distributions on \mathcal{X} , with $\text{supp}(P) \subset \text{supp}(Q)$. Which of the following are true? ($W \circ P$ denotes the output distribution when the input distribution is P and the channel is W .)

- $D(W \circ P \parallel W \circ Q) \leq D(P \parallel Q)$
- $D(W \circ P \parallel W \circ Q) = D(P \parallel Q)$
- $D(W \circ P \parallel W \circ Q) = 0$
- $D(W \circ P \parallel W \circ Q)$ might be ∞

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$D(W \circ P \parallel W \circ Q) \leq D(P \parallel Q)$$

$$D(W \circ P \parallel W \circ Q) = D(P \parallel Q)$$

3) Let X and Y be random variables taking values in alphabet \mathcal{X} . Let P_{XY} be their joint distribution (P_{XY} is a distribution on $\mathcal{X} \times \mathcal{X}$). Which of the **1 point** following are true? ($h(p) = -p \log p - (1-p) \log(1-p)$.)

- $H(X|Y) \leq P_{XY}(X \neq Y) \log |\mathcal{X}| + h(P_{XY}(X \neq Y))$
- $H(Y|X) \leq P_{XY}(X \neq Y) \log |\mathcal{X}| + h(P_{XY}(X \neq Y))$
- $I(X \wedge Y) \geq H(X) - P_{XY}(X \neq Y) \log |\mathcal{X}| - h(P_{XY}(X \neq Y))$
- $I(X \wedge Y) \geq H(Y) - P_{XY}(X \neq Y) \log |\mathcal{X}| - h(P_{XY}(X \neq Y))$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$H(X|Y) \leq P_{XY}(X \neq Y) \log |\mathcal{X}| + h(P_{XY}(X \neq Y))$$

$$H(Y|X) \leq P_{XY}(X \neq Y) \log |\mathcal{X}| + h(P_{XY}(X \neq Y))$$

$$I(X \wedge Y) \geq H(X) - P_{XY}(X \neq Y) \log |\mathcal{X}| - h(P_{XY}(X \neq Y))$$

$$I(X \wedge Y) \geq H(Y) - P_{XY}(X \neq Y) \log |\mathcal{X}| - h(P_{XY}(X \neq Y))$$

4) Let X be a random variable taking values in a finite alphabet $\mathcal{X} \subset \mathbb{R}$. Let P, Q be distributions on \mathcal{X} , such that $D(P \parallel Q) \leq 2$. It is given that **1 point** $\mathbb{E}_Q[X^2] \leq 4$. What can we infer about $\mathbb{E}_P[\log |X|]$?

- $\mathbb{E}_P[\log |X|] \leq 4$
- $\mathbb{E}_P[\log |X|] \leq 2$
- $\mathbb{E}_P[\log |X|] \geq 6$
- $\mathbb{E}_P[\log |X|] \geq 3$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$\mathbb{E}_P[\log |X|] \leq 4$$

$$\mathbb{E}_P[\log |X|] \leq 2$$

5) Let $\mathcal{X} = \{1, 2, 3\}$. Let $W: \mathcal{X} \rightarrow \mathcal{Y}$ be a channel, and let Q be a distribution on \mathcal{Y} . Define $\alpha = \max_P \mathbb{E}_P[D(W_X \parallel Q)]$, where \max is over all **1 point** probability distributions on \mathcal{X} . Let \mathcal{D} be a collection of probability distributions on \mathcal{X} defined as $\mathcal{D} = \{(1, 0, 0), (0, 1, 0), (0, 0, 1)\}$. Which of the following are true?

- $\alpha < \max_{x \in \mathcal{X}} D(W_x \parallel Q)$
- $\alpha = \max_{x \in \mathcal{X}} D(W_x \parallel Q)$
- $\alpha \geq \max_{P \in \mathcal{D}} \mathbb{E}_P[D(W_X \parallel Q)]$
- $\alpha = \max_{P \in \mathcal{D}} \mathbb{E}_P[D(W_X \parallel Q)]$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$\alpha = \max_{x \in \mathcal{X}} D(W_x \parallel Q)$$

$$\alpha \geq \max_{P \in \mathcal{D}} \mathbb{E}_P[D(W_X \parallel Q)]$$

$$\alpha = \max_{P \in \mathcal{D}} \mathbb{E}_P[D(W_X \parallel Q)]$$

6) Let $\mathcal{X} = \{1, 2, 3\}$, $\mathcal{Y} = \{1, 2, 3, 4\}$. Let $W: \mathcal{X} \rightarrow \mathcal{Y}$ be a channel defined as follows: $W_1 = (\frac{1}{2}, \frac{1}{2}, 0, 0)$, $W_2 = (0, 0, 1, 0)$, $W_3 = (0, 0, 0, 1)$. **1 point** Let P be the uniform distribution on \mathcal{X} . Which of the following are true? ($C(W)$ denotes the capacity of channel W , defined as $C(W) = \max_R I(R; W)$, where \max is over all probability distributions on \mathcal{X} .)

- $I(P; W) = H(P) - \mathbb{E}_P[H(W_X)] = \log 3 - \frac{1}{3}$
- $I(P; W) = H(\mathbb{E}_P[W_X]) - \mathbb{E}_P[H(W_X)] = \log 3$
- $C(W) \geq \log 3 - \frac{1}{3}$
- $C(W) \geq \log 3$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$I(P; W) = H(\mathbb{E}_P[W_X]) - \mathbb{E}_P[H(W_X)] = \log 3$$

$$C(W) \geq \log 3 - \frac{1}{3}$$

$$C(W) \geq \log 3$$

7) Let $\mathcal{X} = \mathcal{Y} = \{1, 2, 3\}$. Define $f: \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ as follows: $f(1, 1) = f(2, 2) = f(3, 3) = 1$; $f(i, j) = 0$ for $i \neq j$. Define $g: \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ as **1 point** follows: $g(1, 1) = g(2, 2) = g(3, 3) = 1$; $g(2, 1) = g(2, 3) = 1$; $g(1, 2) = g(1, 3) = g(3, 1) = g(3, 2) = 0$. Let $h: \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ be an arbitrary function. Which of the following are true?

- $\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} f(x, y) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} f(x, y)$
- $\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} g(x, y) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} g(x, y)$
- $\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} h(x, y) \geq \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} h(x, y)$
- $\min_{x \in \mathcal{X}} \min_{y \in \mathcal{Y}} h(x, y) = \min_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} h(x, y)$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} g(x, y) = \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} g(x, y)$$

$$\min_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}} h(x, y) \geq \max_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} h(x, y)$$

$$\min_{x \in \mathcal{X}} \min_{y \in \mathcal{Y}} h(x, y) = \min_{y \in \mathcal{Y}} \min_{x \in \mathcal{X}} h(x, y)$$

8) Let P and Q be distributions on \mathcal{X} , and let $A \subset \mathcal{X}$. Define $f: \mathcal{X} \rightarrow \mathbb{R}$ as $f(x) = I_A(x) - Q(A)$, where $I_A(x) = 1$ if $x \in A$, and **1 point** $I_A(x) = 0$ otherwise. Which of the following are true?

- $\mathbb{E}_P[f(X)] = P(A) - Q(A)$
- $\mathbb{E}_Q[f(X)] = 0$
- $\mathbb{E}_P[f(X)] \leq d(P, Q)$
- $\mathbb{E}_Q[f(X)] \leq d(P, Q)$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$\mathbb{E}_P[f(X)] = P(A) - Q(A)$$

$$\mathbb{E}_Q[f(X)] = 0$$

$$\mathbb{E}_P[f(X)] \leq d(P, Q)$$

$$\mathbb{E}_Q[f(X)] \leq d(P, Q)$$

9) Let $\mathcal{X} = \{0, 1\}$. Let P, Q be distributions on \mathcal{X} , where $P = Q = (0.5 - \delta, 0.5 + \delta)$. Let $Z = \{(0, 0), (0, 1), (1, 0), (1, 1)\}$ and let R be a **1 point** distribution on Z . Which of the following are true?

- $R = (0.25 - \delta, 0.25, 0.25, 0.25 - \delta)$ is a coupling of P and Q .
- $R = (0.25, 0.25 - \delta, 0.25 - \delta, 0.25)$ is a coupling of P and Q .
- $R = (0.25, 0.25 - \delta, 0.25 - \delta, 0.25 + 2\delta)$ is a coupling of P and Q .
- $R = ((0.5 - \delta)^2, 0.25 - \delta^2, 0.25 - \delta^2, (0.5 + \delta)^2)$ is a coupling of P and Q .

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$R = (0.25 - \delta, 0.25, 0.25, 0.25 - \delta) \text{ is a coupling of } P \text{ and } Q.$$

$$R = (0.25, 0.25 - \delta, 0.25 - \delta, 0.25 + 2\delta) \text{ is a coupling of } P \text{ and } Q.$$

$$R = ((0.5 - \delta)^2, 0.25 - \delta^2, 0.25 - \delta^2, (0.5 + \delta)^2) \text{ is a coupling of } P \text{ and } Q.$$

10) Let X be a random variable taking values in \mathcal{X} . Let $f: \mathcal{X} \rightarrow \{0, 1\}^\ell$ and $\phi: \{0, 1\}^\ell \rightarrow \mathcal{X}$ be deterministic functions, such that $\mathbb{P}(\hat{X} \neq X) \leq \epsilon$, where **1 point** $\hat{X} = \phi(f(X))$. Which of the following are true?

- $H(\hat{X}|X) \geq \epsilon \log |\mathcal{X}| + 1$
- $H(\hat{X}|X) \leq \epsilon \log |\mathcal{X}| + 1$
- $H(\hat{X}|X) \geq \ell$
- $I(X \wedge \hat{X}) \leq \ell$

No, the answer is incorrect.

Score: 0

Accepted Answers:

$$H(\hat{X}|X) \leq \epsilon \log |\mathcal{X}| + 1$$

$$I(X \wedge \hat{X}) \leq \ell$$