

X

NPTEL

reviewer2@nptel.iitm.ac.in ▼

Courses » Introduction To Cryptology

Announcements

Course

Ask a Question

Progress



## Unit 5 - Week 4

### Course outline

How to access the portal?

Week 1

Week 2

Week 3

Week 4

- Lecture 1: Cryptographic Hash Functions
- Lecture 2: Random Oracle Model
- Lecture 3: Randomized Algorithm
- Lecture 4: Iterated Construction of Hash Functions
- Lecture 5: Problem Discussions
- Quiz : Week4\_Assignment1
- Feedback form for Week-4
- Assignment Solution

### Week4\_Assignment1

The due date for submitting this assignment has passed. **Due on 2017-08-23, 23:59 IST**  
As per our records you have not submitted this assignment.

1) Let  $(X, Y, K, H)$  be the hash family with  $|X| = 2^{256}$  and  $|Y| = 2^{32}$ . Then the number of all possible hash functions in this family is **1 point**

- $2^{240}$
- $2^{261}$
- $2^{261}$
- $2^{40}$

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

$2^{261}$

2) Assume random oracle model. Suppose that  $h \in F^{(X, Y)}$  is chosen randomly, and let  $X_0 \subseteq X$ . **1 point**  
Suppose that the values  $h(x)$  have been determined (by querying the oracle for  $h$ ) if and only if  $x \in X_0$ . Let  $|X| = N$  and  $|Y| = M$ ,  $N \geq 2M$ . Then choose the correct statement.

- $\Pr[h(x)=y]=1/M$  for all  $x \in X \setminus X_0$  and all  $y \in Y$ .
- $\Pr[h(x)=y]=1/N$  for all  $x \in X \setminus X_0$  and all  $y \in Y$ .
- $\Pr[h(x)=y]=1/(M-|X_0|)$  for all  $x \in X \setminus X_0$  and all  $y \in Y$ .
- $\Pr[h(x)=y]=1/(N-|X_0|)$  for all  $x \in X \setminus X_0$  and all  $y \in Y$ .

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

$\Pr[h(x)=y]=1/M$  for all  $x \in X \setminus X_0$  and all  $y \in Y$ .

3) Let  $(X, Y, K, H)$  be a hash family with  $|Y| = 4096$  and  $X_0 \subset X$  such that  $|X_0| = 32$ . Suppose that  $\epsilon$  is the average-case success probability for finding preimage. Then the best estimate of  $\epsilon$  is **1 point**

- $2^{-12}$
- $2^{-17}$
- $2^{-5}$
- $2^{-7}$

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

$2^{-7}$

4) Let a compression hash function be collision resistant. Then the hash function constructed by Merkle-Damgård algorithm **1 point**

- is collision resistant.
- is not collision resistant.
- may or may not be collision resistant.
- none of them.

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*is collision resistant.*

5) Suppose that  $n = m > 1$  and  $h : Z_{2^m} \rightarrow Z_{2^m}$  is defined by  $h(x) = x^2 + ax + b \pmod{2^m}$ . Then second preimage **1 point**

- can be found only by solving a quadratic equation.
- cannot be found.
- can be found by without solving a quadratic equation.
- Sometimes can be found by solving a linear equation but not always.

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*can be found by without solving a quadratic equation.*

6) Suppose that  $h: X \rightarrow Y$  is a hash function such that it is possible to find  $x, x' \in X$  with  $x \neq x'$  such that  $h(x) = h(x')$ . Then **1 point**

- $h$  is not preimage resistant.
- $h$  is not second preimage resistant.
- $h$  is not collision resistant but may or may not be second preimage resistant.
- $h$  is not collision resistant and not second preimage resistant.

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*$h$  is not collision resistant but may or may not be second preimage resistant.*

7) Find the correct statement among the following. **1 point**

- If a hash function is collision resistant then it is preimage resistant.
- If a hash function is second preimage resistant then it is collision resistant.
- If a hash function is collision resistant then it is second preimage resistant.
- If a hash function is preimage resistant then it is second preimage resistant.

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*If a hash function is collision resistant then it is second preimage resistant.*

8) Suppose that  $h: X \rightarrow Y$  is a hash function considered in random oracle model. Suppose that  $Q$  queries are allowed and  $Q$  is small compared to  $M = |Y|$ . Then the best estimate of the average case success probability of find-second-preimage algorithm is **1 point**

- $Q / M$
- $Q / (M-1)$
- $(Q-1) / (M-1)$
- $(Q-1) / M$

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*$(Q-1) / M$*



9) Let  $B = \{0, 1\}$ . Suppose that  $f: B^{100} \rightarrow B^{50}$  is a collision resistant hash function. Define  $h: B^{200} \rightarrow B^{50}$  such that  $h(x) = f(f(x') \parallel f(x''))$  for all  $x \in B^{200}$ , where  $x = x' \parallel x''$ ,  $x', x'' \in B^{100}$ . Then

1 point

- $h$  is not collision resistant.
- $h$  may or may not be collision resistant.
- $h$  is collision resistant.
- none of them.

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*$h$  is collision resistant.*

10) Suppose you have a 40-bit message digest.  $Q$  is the smallest number of valid pairs required to obtain collision with probability 0.5. Which one of the following with the best estimate of  $Q$ ?

1 point

- $2^{20}$
- $2^{40}$
- $2^{80}$
- $2^{10}$

**No, the answer is incorrect.**

**Score: 0**

**Accepted Answers:**

*$2^{20}$*



Previous Page

End

© 2014 NPTEL - Privacy & Terms - Honor Code - FAQs -

A project of



In association with



Funded by

Government of India  
Ministry of Human Resource Development

Powered by

