

An introduction to Information Theory

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Introduction

Basic blocks of communication system

Lecture #1A: Introduction



Books

Textbook:

- James L. Massey, Lecture notes on “Applied Digital Information Theory I”. (http://www.isiweb.ee.ethz.ch/archive/massey_scr/)



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- Emmanuel Desurvire, "Classical and Quantum Information Theory", 1st Edition, Cambridge University Press, 2009.



Topics to be covered

Information Theory:

- Introduction: Entropy, Relative Entropy, Mutual Information



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- Block to block length coding: Typical sequences



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- Block to block length coding: Typical sequences
- Asymptotic Equipartition Property



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- Coding for sources with memory



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- What are the fundamental limits of data compression?
- Practical source compression algorithm.
- Some mathematical techniques.



Outline of the lecture

- Introduction



Outline of the lecture

- Introduction
- Basic blocks of communication system



Outline

- 1 Introduction
- 2 Basic blocks of communication system



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- Encoding a message at its source.
- Transmitting that message through a communication medium.
- Decoding the message at its destination.



Information Theory

Information Theory answers two fundamental questions in communications

- What is the ultimate data compression?



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- What is the ultimate data compression?
- What is the ultimate transmission rate?



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- A flip of coin with two heads? Does it convey any information?
- A source produces successive bits of $\pi : 3, 1, 4, 1, 5, 9, 2, 6$. Does it convey any information?
- No, no uncertainty in the source.
- Shannon's Information Theory regards only those symbols as information that are not predictable.



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- Main principle of data compression: "Only information essential to understand must be transmitted."



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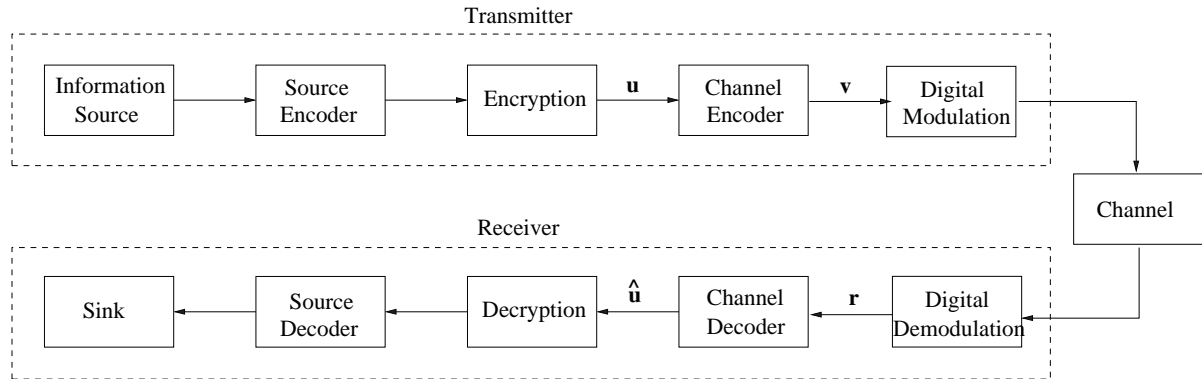


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- Example: If the channel capacity of a particular communication link is (say) 2 Gbps. We can communicate over this channel at any desired rate less than 2 Gbps, and achieve arbitrary low error rates.



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- The output of the source encoder is referred to as the *information sequence*.



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- *Examples:* Data Encryption Standard (DES), RSA system.



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- The encoded sequence that is the output of the channel encoder is referred to as *codeword*.



Modulation

- *Function*: To map the codewords into waveforms which are then transmitted over the physical medium known as the channel.



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- *Examples:* Phase shift keying (PSK), quadrature amplitude modulation (QAM).



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- *Examples:* Binary erasure channel (BEC), Additive white Gaussian noise (AWGN) channel.



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- If the demodulator outputs are unquantized (or has more than two quantization levels), this is known as *soft demodulation*.
- Soft demodulation has significant improvement over hard demodulation.



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- The coded sequences can be broken up into blocks of data *frames*. A frame error occurs if any information bit in that data frame is in error. The decoded FER is the percentage of frames in error.



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- It is in the key that the security of a modern cipher lies, not in the details of the cipher.



Source Decoding

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- *Function:* To reconstruct the original source bits from the decoded information sequence.
- Due to channel errors, the final reconstructed signal may be distorted.

