

### Dynamic Programming Applications

### Water Allocation – Numerical Example

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### **Objectives**

- To demonstrate the water allocation problem through a numerical example using
  - Backward approach
  - Forward approach



### **Numerical Problem**

- Consider a canal supplying water for three different crops
- Maximum capacity of the canal is 4 units of water.



• Optimization Problem: Determine the optimal allocations  $x_i$  to each crop that maximizes the total net benefits from all the three crops



#### Numerical Problem ....contd.

Net benefits from producing the crops can be expressed as a function of the water allotted.  $NB_1(x_1) = 5x_1 - 0.5x_1^2$ 

$$NB_{2}(x_{2}) = 8x_{2} - 1.5x_{2}^{2}$$
$$NB_{3}(x_{3}) = 7x_{3} - x_{3}^{2}$$

The net benefit values are calculated for each crop and are as shown

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	<u>Lable L</u>										
	$x_i$	$NB_1(x_1)$	$NB_2(x_2)$	$NB_3(x_3)$							
	0	0.0	0.0	0.0							
	1	4.5	6.5	6.0							
	2	8.0	10.0	10.0							
	3	10.5	10.5	12.0							
	4	12.0	8.0	12.0							
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#### Numerical Problem ....contd.

**X**<sub>1</sub>

Representation of the problem as a set of nodes and links

**X**3

**X**<sub>2</sub>



**Optimization Methods: M6L4** 



#### Numerical Problem ...contd.

- The values inside the node show the value of state variable at each stage
- Number of nodes for any stage corresponds to the number of discrete states possible for each stage.
- The values over the links show the different values taken by decision variables corresponding to the value taken by state variables



# Numerical Problem: Solution by Backward recursion

Sub-optimization function for the 3<sup>rd</sup> crop:

 $f_3(S_3) = \max_{\substack{x_3 \ 0 \le x_3 \le S_3}} NB_3(x_3)$  with the range of  $S_3$  from 0 to 4.

Table 2

State		f.(S.)	*					
.S <sub>3</sub>	<i>x</i> <sub>3</sub> :	0	1	2	3	4	- /3(03)	×3
0		0				•	0	0
1	·	0	б				6	1
2		0	б	10			10	2
3		0	б	10	12		12	3
4	•	0	б	10	12	12	12	3,4

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• Next, by considering last two stages together, the sub-optimization function is  $f_2(S_2) = \max_{\substack{x_2 \\ x_2 \le S_2}} [NB_2(x_2) + f_1(S_2 - x_2)]$ 

Table 3

The calculations are shown below

State S <sub>2</sub>	<i>x</i> <sub>2</sub>	$NB_2(x_2)$	$(S_2 - x_2)$	$f_3(S_2 - x_2)$	$\begin{array}{l} f_2(S_2) = \\ NB_2(x_2) + \\ f_3(S_2 - x_2) \end{array}$	$f_2^{*}(S_2)$	x2*
0	0	0	0	0	0	0	0
1	0	0	1	6	6	6.5	1
I _	1	6.5	0	0	6.5	0.5	1
Table con	ntd. on n	ext slide				• •	



	0	0	2	10	10	•	
2 -	1	6.5	1	6	12.5	12.5	1
-	2	10	0	0	10	_	
	0	0	3	12	12	• • • • •	
2	1	6.5	2	10	16.5	165	1
5 -	2	10	1	6	16	_ 10.5	1
-	3	10.5	0	0	10.5	_	
	0	0	4	12	12		
-	1	6.5	3	12	18.5	_	
4 -	2	10	2	10	20	20	2
-	3	10.5	1	6	16.5	_	
	4	8	0	0	8	_	



Considering all the three stages together

$f_1(Q) = \max[NB_1(x_1) + f_2(Q - x_1)]$	with	$S_1 = Q = 4$
$x_1 \\ 0 \le x_1 \le Q$		

State $S_1 = Q$	x <sub>1</sub>	$NB_1(x_1)$	$(Q-x_1)$	$f_2(Q-x_1)$	$f_1(S_1) =$ $NB_1(x_1) +$ $f_2(Q - x_1)$	$f_1^*(S_1)$	x1*
	0	0	4	20	20	•	
	1	4.5	3	16.5	21	-	
4	2	8	2	12.5	20.5	21	1
	3	10.5	1	6.5	17	-	
	4	12	0	0	12	-	

Table 4



- Backtrack through each table,
- Optimal allocation for crop 1,  $x_1^* = 1$  and  $S_1 = 4$
- Thus,  $S_2 = S_1 x_1 = 3$
- From 2<sup>nd</sup> stage, the optimal allocation for crop 2,  $x_2 = 1$ .
- Now,  $S_3 = S_2 x_2 = 2$
- From 3<sup>rd</sup> stage calculations,  $x_3^* = 2$
- Maximum total net benefit from all the crops = 21



# Numerical Problem: Solution by Forward recursion

- Start from the first stage and proceed towards the final stage
- Suboptimization function for the first stage

$$f_1(S_1) = \max_{x_1 \le S_1} NB_1(x_1)$$

- Range of values for  $S_1$  is from 0 to 4
- The calculations are shown in the table



		<u>Table 5</u>		
State S <sub>1</sub>	<i>x</i> <sub>1</sub>	$NB_1(x_1)$	$f_2^{*}(S_2)$	x1*
0	0	0	0	0
1	0	0		. 1
1	1	4.5	- 4.5	1
	0	0	•	•
2	1	4.5	8	2
	2	8		
	0	0	- _ 10.5 -	•
2	1	4.5		3
5	2	8		
	3	10.5		
	0	0		
	1	4.5	-	
4	2	8	12	4
	3	10.5	-	
	4	12	-	
	•	•	•	•

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## Solution by Forward recursion ...

ooman	State S <sub>2</sub>	<i>x</i> <sub>2</sub>	$NB_2(x_2)$	$(S_2 - x_2)$	$f_1(S_2 - x_2)$	$\begin{array}{l} f_{2}(S_{2}) = \\ NB_{2}(x_{2}) + \\ f_{2}(S_{2} - x_{2}) \end{array}$	$f_2^{*}(S_2)$	x2*
$f_2(S_2) =$	0	0	0	0	0	0	0	0
$\begin{bmatrix} \mathbf{N}\mathbf{D} & (\mathbf{n} \mathbf{n} \mathbf{n} \mathbf{n} \mathbf{n} \end{bmatrix}$	1	0	0	1	4.5	4.5	6.5	1
$\max   NB_2(x_2) +  $	1 _	1	6.5	0	0	6.5	0.5	1
$\frac{x_2}{x_2} \int_{-\infty} f_1(S_2 - x_2)$		0	0	2	8	8	· · ·	
$x_2 \leq S_2 \square \circ 1 + 2 = 2 + \square$	2 -	1	6.5	1	4.5	11	11	1
	-	2	10	0	0	10	-	
		0	0	3	10.5	10.5		
	-	1	6.5	2	8	14.5	145	1,2
	<u> </u>	2	10	1	4.5	14.5	_ 14.5	
	-	3	10.5	0	0	10.5		
		0	0	4	12	12		
	-	1	6.5	3	10.5	17	-	
	4	2	10	2	8	18	18	2
	-	3	10.5	1	4.5	15	-	
	-	4	8	0	0	8	-	
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• 
$$f_3(S_3) = \max_{\substack{x_3 \\ x_3 \le S_3 = Q}} \left[ NB_3(x_3) + f_2(S_3 - x_3) \right]$$
 with  $S_3 = 4$ 

<u>Table 7</u>										
State S <sub>3</sub>	<i>x</i> <sub>3</sub>	$NB_3(x_3)$	S <sub>3</sub> - x <sub>3</sub>	$f_2(S_3 - x_3)$	$f_3(S_3) = \\NB_3(x_3) + \\f_2(S_3 - x_3)$	$f_{3}^{*}(S_{3})$	x3*			
•	0	0	4	18	18	· · ·				
	1	б	3	14.5	20.5					
4	2	10	2	. 11	21	21	2			
	3	12	1	6.5	18.5					
	4	12	0	0	12					



- Backtrack through each table,
- Optimal allocation for crop 3,  $x_3^* = 2$  and  $S_3 = 4$
- Then,  $S_2 = S_3 x_3 = 2$
- The optimal allocation for crop 2,  $x_2^* = 1$

• Now, 
$$S_1 = S_2 - x_2 = 1$$

- From 1<sup>st</sup> stage calculations,  $x_1^* = 1$
- Maximum total net benefit from all the crops = 21
- These solutions are the same as those we got from backward recursion method.



## Thank You

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